

TREE ORDINALS AND PROOF- THEORETIC BOUNDING FUNCTIONS

OR

"What I learned from WB".

Stan Wainer (Leeds)

Buchholz papers that influenced me

1980 3 Contribs. "Recent Advances in Proof Thy"
Oxford.

1981 + Fef. Pohlers & Sieg "Iterated Ind.
Defns & Subsystems of Anal" LNM 897

1987 "Indep. Result for $(\Pi_1^1 CA) + (BI)$ " APAL

1994 + Cichon & Weiermann "Unif. Approach
to fund. seqs & subrec. hierarchies" MLQ

Ordinal Presentations

$$\alpha = \bigcup_n \alpha[0] \subset \alpha[1] \subset \dots \subset \alpha[n] \subset \dots$$

where $\alpha[n]$ finite

$$\beta+1 \in \alpha[n] \Rightarrow \beta \in \alpha[n]$$

$$\beta \in \alpha[n] \Rightarrow \beta+1 \in \alpha[n+1]$$

(B-C-W : β has α -norm $\leq n$)

Predecessors : $P_n(\alpha) = \max \alpha[n]$

Rank Function : $G_\alpha(n) = \text{size } \alpha[n]$

Fast-Growing : $B_\alpha(n) = B_{P_n(\alpha)}^2(n)$

Tree Ordinals (Ω_1)

- $0 \in \Omega_1$ $0[n] = \emptyset$
- $\alpha \in \Omega_1 \Rightarrow \alpha+1 \in \Omega_1$ $\alpha+1[n] = \alpha[n] \cup \{\alpha\}$
- $\lambda: \mathbb{N} \rightarrow \Omega_1 \Rightarrow \lambda \in \Omega_1$ $\lambda[n] = \lambda_n[n]$
 $\lambda_i \in \lambda_{i+1}[i+1]$

Then $\alpha[n]$ is a presentation of $\|\alpha\|$

(+) $\beta+0=\beta$, $\beta+(\alpha+1)=(\beta+\alpha)+1$, $\beta+\lambda = \langle \beta+\lambda_i \rangle_{i \in \mathbb{N}}$
 (exp) $2^0=1$, $2^{\alpha+1} = 2^\alpha + 2^\alpha$, $2^\lambda = \langle 2^{\lambda_i} \rangle_{i \in \mathbb{N}}$.

Tree Ordinals (Ω_2)

- $0 \in \Omega_2$ $0[\delta, n] = \emptyset$
- $\alpha \in \Omega_2 \Rightarrow \alpha+1 \in \Omega_2$ $\alpha+1[\delta, n] = \alpha[\delta, n] \cup \{\alpha\}$
- $\lambda: \mathbb{N} \rightarrow \Omega_2 \Rightarrow \lambda \in \Omega_2$ $\lambda[\delta, n] = \lambda_n[\delta, n]$
- $\lambda: \Omega_1 \rightarrow \Omega_2 \Rightarrow \lambda \in \Omega_2$ $\lambda[\delta, n] = \lambda_\delta[\delta, n]$
 $\xi \in \delta[n] \Rightarrow \lambda_\xi \in \lambda_\delta[\delta, n]$

Fast-Growing $\varphi: \Omega_2 \times \Omega_1 \rightarrow \Omega_1$

$$\left[\begin{array}{ll} \varphi_0(\beta) = \beta + 1 & \\ \varphi_{\alpha+1}(\beta) = \varphi_\alpha \circ \varphi_\alpha(\beta) & \\ \varphi_\lambda(\beta) = \langle \varphi_{\lambda_i}(\beta) \rangle_{i \in \mathbb{N}} & \lambda: \mathbb{N} \rightarrow \Omega_2 \\ \varphi_\lambda(\beta) = \varphi_{\lambda\beta}(\beta) & \lambda: \Omega_1 \rightarrow \Omega_2 \end{array} \right.$$

Then with $\omega = \text{id}_\mathbb{N}$ and $\omega_1 = \text{id}_{\Omega_1}$:

$$\varphi_{\omega_1}(\beta) = \beta + 2^\beta$$

$$\varphi_{\omega_1 \cdot 2}(\omega) = \varphi_{\omega_1}^{2^\omega}(\omega) = \dot{\epsilon}_0$$

$$\varphi_{\dot{\epsilon}_{\omega_1+1}}(\omega) = \text{Howard}$$

G-Collapse: $G_{\varphi_\alpha(\beta)}^{(n)} = B_{G_\alpha(n)}(G_\beta(n))$

Infinitary Arithmetic $n: N \vdash^\alpha \Gamma$

• Logic :-

(Ax) $n: N \vdash^\alpha \Gamma$ if a true atom $\in \Gamma$

($\exists$)
$$\frac{n: N \vdash^\beta m: N \quad n: N \vdash^\beta \Gamma, A(m)}{\beta \in \alpha[n] \quad n: N \vdash^\alpha \Gamma, \exists x A(x)}$$

($\forall$)
$$\frac{\max(n, i): N \vdash^{\beta_i} \Gamma, A(i) \text{ all } i}{\beta_i \in \alpha[\max n, i] \quad n: N \vdash^\alpha \Gamma, \forall x A(x)}$$

+ Propositional rules + Cut.

• Computation :-

(N1) $n: N \vdash^\alpha \Gamma, m: N$ if $m \leq n+1$

(N2)
$$\frac{n: N \vdash^\beta m: N \quad m: N \vdash^\beta \Gamma}{n: N \vdash^\alpha \Gamma}$$

Embedding of PA

$$PA \vdash_r^k A(x) \Rightarrow n:N \vdash_r^{\omega \cdot k} A(n)$$

Proof

Suppose $\forall x A(x)$ proven by Induction;
from $\vdash^{k-1} A(0)$ and $\vdash^{k-1} \forall x (A(x) \rightarrow A(x+1))$

$$\text{Cut } \frac{A(0) \quad A(0) \rightarrow A(1)}{A(1)}$$

$$\text{Cut } \frac{A(1) \quad A(1) \rightarrow A(2)}{A(2)}$$

$$\text{Cut } \frac{A(2) \quad A(2) \rightarrow A(3)}{A(3)}$$

... etcetera

$\therefore$ by successive A-Cuts :-

$$n:N \vdash_r^{\omega \cdot (k-1) + n} A(n)$$

$$\therefore \text{By } (\forall)^\omega \vdash_r^{\omega \cdot (k-1) + \omega} \forall x A(x)$$

Cut Elimination (Gentzen, Schütte, ...)

1. $n \vdash_r^\beta \neg C, B ; n \vdash_r^\delta C, B \Rightarrow n \vdash_r^{\beta+\delta} B$
provided $\delta \in \mathbb{I} \subseteq \beta \in \mathbb{I}$ and $|C| \leq r+1$.

2. Hence $n \vdash_{r+1}^\alpha B \Rightarrow n \vdash_r^{2^\alpha} B$.

Proof by induction on δ .

Suppose $C \equiv \exists x D(x)$, and

$\delta \in \delta[n]$:
$$\frac{n \vdash^\delta m \quad n \vdash^\delta C, D(m), B}{n \vdash^\delta C, B}$$

$\forall$ -Invert $n \vdash^\beta \neg C, B$ to get $m \vdash^\beta \neg D(m), B$.

Then by the ind. hypoth. at $\delta < \delta$,

(N₂)
$$\frac{n \vdash^\delta m \quad m \vdash^\beta \neg D(m), B}{n \vdash^{\delta+\beta} \neg D(m), B}$$

D-Cut
$$\frac{n \vdash^{\beta+1} \neg D(m), B \quad n \vdash^{\beta+\delta} D(m), B}{n \vdash_r^{\beta+\delta} B}$$

Bounding Functions B_α

$$(\exists) \frac{n: N \vdash_\beta m: N \quad n: N \vdash \text{Spec}(n, m)}{n: N \vdash_\alpha \exists y \text{Spec}(n, y)}$$

$$\text{Let } B_\beta(n) = \max \{m \mid n: N \vdash_\beta m: N\}$$

$$\text{If } \frac{n: N \vdash_\gamma k: N \quad k: N \vdash_\gamma m: N}{n: N \vdash m: N}$$

$$\text{then } m \leq B_\gamma(k) \wedge k \leq B_\gamma(n)$$

$$\text{so } B_{\gamma+1}(n) = B_\gamma \circ B_\gamma(n).$$

To complete the definition add

$$B_0(n) = n+1, \quad B_\lambda(n) = B_{\lambda_n}(n).$$

Theorem If $\text{PA} \vdash \forall x \exists y \text{Spec}(x, y)$

$\forall n \exists m < B_\alpha(n) \text{Spec}(n, m)$
for some $\alpha < \varepsilon_0$.

Fast-Growing Hierarchy

$$B_0(n) = n+1$$

$$B_{\alpha+1}(n) = B_\alpha \circ B_\alpha(n)$$

$$B_\lambda(n) = B_{\lambda_n}(n)$$

Can be written as

$$B_\alpha(n) = n \oplus 2^\alpha$$

$$n \oplus 2^0 = n+1$$

$$n \oplus 2^{\alpha+1} = n \oplus 2^\alpha \oplus 2^\alpha$$

$$n \oplus 2^\lambda = n \oplus 2^{\lambda_n}$$

Infinitary ID_1 $\delta: \Omega_1, n: N \vdash^{\alpha \in \Omega_2} \Gamma$

$$ID_1(W)^\infty = \text{Inf. Arith} + (W) + (\Omega_1)$$

$$(W) \quad \frac{\delta, n \vdash^\beta m: N \quad \delta, n \vdash^\beta \Gamma, Ol(W, m)}{\delta, n \vdash^\alpha \Gamma, W(m)}$$

where $\beta \in \alpha[\delta, n]$ and $Ol(W, m)$ is
 $m = 0 \vee \exists k \in W (m = k \oplus 1) \vee \forall i \exists k \in W (\{m\}(i) \approx k)$

$$(Q) \quad \frac{\begin{array}{c} \omega, l \vdash_0^\delta \Delta, W(m) \\ \delta \in \Omega_1, \delta \in \delta[l], \Delta \subset Pos W \end{array} \quad \Downarrow \quad \begin{array}{c} \delta, n \vdash^\beta \Gamma_0, W(m) \quad \delta, \max(n, l) \vdash^\beta \Delta, \Gamma_1 \end{array}}{\delta, n \vdash^\alpha \Gamma_0, \Gamma_1}$$

$\beta \in \alpha[\delta, n]$.

Embedding $ID_1(W)$

$$ID_1 \vdash \Gamma(\vec{x}) \Rightarrow \omega, n \vdash^{\omega_1 + \omega} \Gamma(\vec{n})$$

Proof To derive the W axioms :-

$$(W1) \quad \forall x (\mathcal{O}(W, x) \rightarrow W(x))$$

Straightforward by W-Rule and $\forall^\omega$

$$(W2) \quad \forall x (\mathcal{O}(A, x) \rightarrow A(x)) \rightarrow W(m) \rightarrow A(m)$$

Apply Ω -Rule with $\beta = \omega$, and

$$\Gamma_0 = \overline{W(m)}$$

$$\Gamma_1 = \neg \forall x (\mathcal{O}(A, x) \rightarrow A(x)), A(m)$$

The point is that $\omega, l \vdash_0^\delta \Delta, W(m)$
can be transformed into

$$\delta, l \vdash^\delta \Delta, \Gamma_1 \quad \text{hence} \quad \delta, l \vdash^{\omega_1} \Delta, \Gamma_1$$

by weakening since $\omega_1[\delta, l] = \delta[\delta, l]$.

Collapsing

If $\gamma, n \vdash_0^\alpha \Gamma \in \text{Pos } W_1$ in $ID_1(W)^\infty$
then $n \vdash_0^{\varphi_\alpha(\gamma)} \Gamma$ without (Ω_1)

Proof

$$\delta, l \vdash_0^\delta \Delta, W(m)$$

$\Downarrow$

$$\begin{array}{c} (\Omega_1) \quad \gamma, n \vdash_0^\beta \Gamma_0, W(m) \quad \delta, l \vdash_0^\beta \Delta, \Gamma_1 \\ \hline \gamma, n \vdash_0^\alpha \Gamma_0, \Gamma_1 \end{array}$$

$$\text{IHyp (LHP)} \quad n \vdash_0^{\varphi_\beta(\gamma)} \Gamma_0, W(m)$$

$$\text{Apply RHP} \quad \varphi_\beta(\gamma), n \vdash_0^\beta \Gamma_0, \Gamma_1$$

$$\text{IHyp (RHP)} \quad n \vdash_0^{\varphi_\beta \varphi_\beta(\gamma)} \Gamma_0, \Gamma_1$$

Note $\beta \in \alpha[\gamma, n] \Rightarrow \varphi_\beta(\gamma) \in \varphi_\alpha(\gamma)[n]$
 $\Rightarrow \varphi_\beta^2(\gamma)[n] \subseteq \varphi_\alpha(\gamma)[n]$

Bounding

$$ID_1(W) \vdash \exists y \text{Spec}(x, y)$$

$$\omega, n \vdash^{\omega, 2} \exists y \text{Spec}(n, y)$$

$$\omega, n \vdash_0^{\alpha} \exists y \text{Spec}(n, y) \quad \alpha = 2 \cdots 2^{\omega, 2}$$

$$n \vdash_0^{\varphi_{\alpha}(\omega)} \exists y \text{Spec}(n, y)$$

$$\models \exists y < B_{\varphi_{\alpha}(\omega)}(n). \text{Spec}(n, y)$$

Fast-Growing below "Howard".

And so on - $ID_{<\omega}(W) \equiv \Pi_1^1\text{-CA}_0$