

Finite trees as ordinals

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Honouring Wilfried
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Typical trees

The natural numbers:

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 - ▶ Length of sequences
 - ▶ Rightmost element where they differ

Elementary properties

$$\mathbf{A} < \mathbf{B} \Leftrightarrow \mathbf{A} \leq \langle \mathbf{B} \rangle \vee (\langle \mathbf{A} \rangle < \mathbf{B} \wedge \langle \mathbf{A} \rangle < \langle \mathbf{B} \rangle)$$

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- ▶ Decidable
- ▶ Transitive
- ▶ Linear
- ▶ Equality is the usual tree equality

Some ordinal functions

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$$\overset{\alpha}{!} = \alpha + 1$$

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where \sim means we jump over fix points.

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In general we get the fix point free n -ary Veblen functions.

Approximating from below 1

$$\Gamma_0 = \begin{array}{c} \cdot \\ \diagdown \quad \cdot \\ \cdot \quad \diagup \\ \cdot \quad \cdot \quad \cdot \\ \cdot \quad \cdot \quad \cdot \quad \cdot \end{array}$$

Start with immediate subtrees:

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This gives all trees less than Γ_0 .

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This gives all trees less than Γ_0 . To get a cofinal set we only need

$$\begin{array}{c} \gamma \\ \cdot \quad \cdot \\ \cdot \quad \cdot \quad \cdot \end{array}$$

Wellfoundedness

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Both arguments are straightforward.

Further work

Linear extensions of embeddings

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$$\left| \begin{array}{c} A \quad B \\ \diagdown \quad / \\ \cdot \end{array} \right| \leq \left| \begin{array}{c} |A| \quad |B| \\ \diagdown \quad / \\ \cdot \end{array} \right| \oplus \left| \begin{array}{c} |B| \quad |A| \\ \diagdown \quad / \\ \cdot \end{array} \right|$$

This gives Higman's lemma. Further work gives Kruskal's theorem.

Further work

Finite trees with labels

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- ▶ $A <_i B \Leftrightarrow A \leq_i \langle B \rangle_i \vee (\langle A \rangle_i < B \vee A <_{i+} B)$

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- ▶ $A <_\infty B$ — lexicographical ordering

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- ▶ Takeuti's ordinal diagrams